# Geometric properties of random matrices with independent log-concave rows/columns 

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Based on joint work with
O. Guédon, A. Litvak, A. Pajor, N. Tomczak-Jaegermann

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- $X$ is $\psi_{\alpha}(\alpha \in[1,2])$ with constant $C$ if for all $y \in \mathbb{R}^{n}$,

$$
\|\langle X, y\rangle\|_{\psi_{\alpha}} \leq C|y|
$$

where

$$
\|Y\|_{\psi_{\alpha}}=\inf \left\{a>0: \mathbb{E} \exp \left((Y / a)^{\alpha}\right) \leq 2\right\}
$$

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If for a set $K \subseteq \mathbb{R}^{n}$ the random vector distributed uniformly on $K$ is isotropic, we say that $K$ is isotropic.

## Log-concavity

A random vector $X$ in $\mathbb{R}^{n}$ is log-concave if its law $\mu$ satisfies a Brunn-Minkowski type inequality

$$
\mu(\theta A+(1-\theta) B) \geq \mu(A)^{\theta} \mu(B)^{1-\theta}
$$

## Theorem (Borell)

A random vector not supported on any $(n-1)$ dimensional hyperplane is log-concave iff it has density of the form $\exp (-V(x))$, where $V: \mathbb{R}^{n} \rightarrow(-\infty, \infty]$ is convex.

## Lemma (Borell)

An isotropic log-concave random vector is $\psi_{1}$ with a universal constant $C$.

## Examples

The following distributions are log-concave:

- Gaussian measures
- Uniform distributions on convex bodies
- Measures with density of the form $C \exp (-\|x\|)$, where $\|x\|$ is a norm.
- Products, affine images and convolutions of the above distributions.


## The basic model

## Definition

Let $\Gamma$ be an $n \times N$ matrix with columns $X_{1}, \ldots, X_{N}$, where $X_{i}$ 's are independent isotropic log-concave random vectors in $\mathbb{R}^{n}$

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- What is the operator norm of $\Gamma: \ell_{2}^{N} \rightarrow \ell_{2}^{n}$ ?
- When is $\Gamma^{T}$ close to a multiple of isometry?
- How does $\Gamma$ act on sparse vectors?
- What is the smallest singular value of $\Gamma$ ?


## Motivations: sampling convex bodies

## Problem

Let $K \subseteq \mathbb{R}^{n}$ be a convex body, s.t. $B_{2}^{n} \subseteq K \subseteq R B_{2}^{n}$. Assume we have access to an oracle (a black box), which given $x \in \mathbb{R}^{n}$ tells us whether $x \in K$.

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How to generate random points uniformly distributed in $K$ ?
How to compute the volume of $K$ ?

- This can be done by using Markov chains.
- Their speed of convergence depends on the position of the convex body.
- Preprocessing: First put $K$ in the isotropic position (again by randomized algorithms).
- Centering the body is not comp. difficult - takes $\mathcal{O}(n)$ steps.
- The question boils down to:

How to approximate the covariance matrix of $X$ - uniformly distributed on $K$ by the empirical covariance matrix

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or (after a linear transformation)
Given an isotropic convex body in $\mathbb{R}^{n}$, how large $N$ should we take so that

$$
\left\|\frac{1}{N} \sum_{i=1}^{N} X_{i} \otimes X_{i}-I d\right\|_{\ell_{2} \rightarrow \ell_{2}} \leq \varepsilon
$$

with high probability?

## Interpretation in terms of $\Gamma$.

## We have

$$
\begin{aligned}
\left\|\frac{1}{N} \sum_{i=1}^{N} X_{i} \otimes X_{i}-I d\right\|_{\ell_{2} \rightarrow \ell_{2}} & =\sup _{y \in S^{n-1}}\left|\frac{1}{N} \sum_{i=1}^{N}\left\langle X_{i}, y\right\rangle^{2}-1\right| \\
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Let $\Gamma$ be a matrix with independent columns $X_{1}, \ldots, X_{N}$ drawn from an isotropic convex body (log-concave measure) in $\mathbb{R}^{n}$.

How large should $N$ be so that $N^{-1 / 2} \Gamma^{T}: \mathbb{R}^{n} \rightarrow \mathbb{R}^{N}$ was an almost isometry?

## History of the problem

- Kannan, Lovasz, Simonovits (1995) $-N=\mathcal{O}\left(n^{2}\right)$


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For arbitrary isotropic random vectors, if you do not assume any uniform bound on $\left\langle X_{i}, y\right\rangle, y \in S^{n-1}$, you cannot remove the logarithm (the optimal bound $N=\mathcal{O}\left(n \log ^{\beta} n\right)$ is due to M . Rudelson). Recently $N=O(n \log \log n)$ was proven under a uniform bound on $(4+\varepsilon)$-th moments of $\left\langle X_{i}, y\right\rangle$ (R. Vershynin).

## Remark

If $\frac{1}{\sqrt{N}} \Gamma^{T}$ is an almost isometry then obviously $\|\Gamma\| \leq C \sqrt{N}$, so the KLS question and the question about $\|\Gamma\|$ are related.

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It turns out that to answer KLS it is enough to have good bounds on

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A_{m}:=\sup _{\substack{z \in S^{N-1} \\|\operatorname{supp} z| \leq m}}|\Gamma z|
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## Theorem (Litvak, Pajor, Tomczak-Jaegermann, R.A.)

If $N \leq \exp (c \sqrt{n})$ and the vectors $X_{i}$ are log-concave then for $t>1$, with probability at least $1-\exp (-c t \sqrt{n})$,

$$
\forall_{m \leq N} A_{m} \leq C t\left(\sqrt{n}+\sqrt{m} \log \left(\frac{2 N}{m}\right)\right) .
$$

In particular, with high probability $\|\Gamma\| \leq C(\sqrt{n}+\sqrt{N})$.

## Sketch of the proof

A modification of Bourgain's approach. One approximates an arbitrary vector $z$ with $|\operatorname{supp} z| \leq m$ by $x_{0}+x_{1}+\ldots+x_{l}\left(l<\log _{2} m\right)$, where

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\begin{array}{r}
\left|\operatorname{supp} x_{i}\right| \simeq m / 2^{i},\left\|x_{i}\right\|_{\infty} \simeq \sqrt{2^{i} / m}, i \geq 1 \\
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and $x_{i}$ comes from a $2^{-i}$-net in the set of sparse vectors of support at most $m / 2^{i}$.

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Then using the $\psi_{1}$ condition one shows that with high probability

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A_{m}^{2} \lesssim \max _{i}\left|X_{i}\right|^{2}+A_{m}(\sqrt{n}+\sqrt{m} \log (2 N / m))
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Thus max $_{i}\left|X_{i}\right| \leq C \sqrt{n}$ with high probability and we can solve the inequality for $A_{m}$.

## Compressed sensing and neighbourly polytopes

Imagine we have a vector $x \in \mathbb{R}^{N}$ ( $N$ large), which is supported on a small number of coordinates (say $|\operatorname{supp} \mathrm{x}|=m \ll N$ ).

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If we knew the support of $x$, to determine $x$ it would be enough to take $m$ measurements along basis vectors.

What if we don't know the support?
Answer (Donoho, Candes, Tao, Romberg) Take measurements in random directions $Y_{1}, \ldots, Y_{n}$ and set

$$
\hat{x}=\operatorname{argmin}\left\{\|y\|_{1}:\left\langle Y_{i}, y\right\rangle=\left\langle Y_{i}, x\right\rangle\right\}
$$

## Compressed sensing and neighbourly polytopes

## Definition

A polytope $K \subseteq \mathbb{R}^{n}$ is called m-neighbourly if any set of vertices of $K$ of cardinality at most $m+1$ is the vertex set of a face.

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A (centraly symetric) polytope $K \subseteq \mathbb{R}^{n}$ is called
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## Theorem (Donoho)

Let $\Gamma$ be an $n \times N$ matrix with columns $X_{1}, \ldots, X_{N}$. The following conditions are equivalent
(i) For any $x \in \mathbb{R}^{N}$ with $|\operatorname{supp} x| \leq m, x$ is the unique solution of the minimization problem

$$
\min \|t\|_{1}, \quad \Gamma t=\lceil x .
$$

(ii) The polytope $K(\Gamma)=\operatorname{conv}\left( \pm X_{1}, \ldots, \pm X_{N}\right)$ has $2 N$ vertices and is $m$-symmetric-neighbourly.

## Compressed sensing and neighbourly polytopes

## Definition (Restricted Isometry Property (Candes, Tao))

For an $n \times N$ matrix $\Gamma$ define the isometry constant $\delta_{m}=\delta_{m}(\Gamma)$ as the smallest number such that

$$
\left(1-\delta_{m}\right)|x|^{2} \leq|\Gamma x|^{2} \leq\left(1+\delta_{m}\right)|x|^{2}
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for all $m$-sparse vectors $x \in \mathbb{R}^{N}$.

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## Theorem (Candes)

If $\delta_{2 m}(\Gamma)<\sqrt{2}-1$ then for every $m$-sparse $x \in \mathbb{R}^{n}, x$ is the unique solution to

$$
\min \|t\|_{1}, \quad \Gamma t=\Gamma x
$$

In consequence, the polytope $K(\Gamma)\left(r e s p . K_{+}(\Gamma)=\operatorname{conv}\left(X_{1}, \ldots, X_{N}\right)\right)$ is m-symmetric-neighbourly (resp. m-neighbourly)

## History

The following matrices satisfy RIP

- Gaussian matrices (Candes, Tao), $m \simeq n / \log (2 N / n)$
- Matrices with rows selected randomly from the Fourier matrix (Candes \& Tao, Rudelson \& Vershynin), $m \simeq n / \log ^{4}(N)$
- Matrices with independent subgaussian isotropic rows (Mendelson, Pajor, Tomczak-Jaegermann), $m \simeq n / \log (2 N / n)$
- Matrices with independent log-concave isotropic columns (LPTA), $m \simeq n / \log ^{2}(2 N / n)$


## Neighbourly polytopes

## Theorem (LPTA)

Assume that $X_{i}^{\prime} s$ are $\psi_{r}$. Let $\theta \in(0,1 / 4)$ and assume that $N \leq \exp \left(c \theta^{C} n^{c}\right)$ and $m \log ^{2 / r}\left(\frac{2 N}{\theta m}\right) \leq \theta^{2} n$. Then, with probability at least $1-\exp \left(-c \theta^{c} n^{c}\right)$

$$
\delta_{m}\left(\frac{1}{\sqrt{n}} \Gamma\right) \leq \theta .
$$

## Corollary (LPTA)

Let $X_{1}, \ldots, X_{N}$ be random vectors drawn from an isotropic $\psi_{r}$ ( $r \in[1,2]$ ) convex body in $\mathbb{R}^{n}$. Then, for $N \leq \exp \left(c n^{c}\right)$, with probability at least $1-\exp \left(-c n^{c}\right)$, the polytope $K(\Gamma)\left(r e s p . K_{+}(\Gamma)\right)$ is m-symmetric-neighbourly (resp. m-neighbourly) with

$$
m=\left\lfloor c \frac{n}{\log ^{2 / r}(C N / n)}\right\rfloor
$$

## Method of proof

We use the same approximation techniques as for the KLS problem to bound

$$
B_{m}=\sup _{|\operatorname{supp} z| \leq m,|z|=1} \|\left.\sum_{i \leq N} z_{i} X_{i}\right|^{2}-\left.\sum_{i \leq N} z_{i}^{2}\left|X_{i}\right|^{2}\right|^{1 / 2}
$$

## Theorem (B. Klartag)

$$
\mathbb{P}\left(\max _{i \leq N}\left|\frac{\left|X_{i}\right|^{2}}{n}-1\right| \geq \varepsilon\right) \leq C \exp \left(-c \varepsilon^{c} n^{c}\right)
$$

Thus

$$
\delta_{n}\left(n^{-1 / 2} \Gamma\right) \leq n^{-1} B_{m}^{2}+\varepsilon
$$

with overwhelming probability.

## Smallest singular value

## Definition

For an $n \times n$ matrix $\Gamma$ let $s_{1}(\Gamma) \geq s_{2}(\Gamma) \geq \ldots \geq s_{n}(\Gamma)$ be the singular values of $\Gamma$, i.e. eigenvalues of $\sqrt{\Gamma \Gamma^{\top}}$. In particular

$$
s_{1}(\Gamma)=\|A\|, s_{n}(\Gamma)=\inf _{x \in S^{n-1}}|\Gamma x|=\frac{1}{\left\|A^{-1}\right\|}
$$

## Theorem (Edelman, Szarek)

Let $\Gamma$ be an $n \times n$ random matrix with independent $\mathcal{N}(0,1)$ entries. Let $s_{n}$ denote the smallest singular values of $\Gamma$. Then, for every $\varepsilon>0$,

$$
\mathbb{P}\left(s_{n}(\Gamma) \leq \varepsilon n^{-1 / 2}\right) \leq C \varepsilon
$$

where $C$ is a universal constant.

## Theorem (Rudelson, Vershynin)

Let $\Gamma$ be a random matrix with independent entries $X_{i j}$, satisfying $\mathbb{E} X_{i j}=0, \mathbb{E} X_{i j}^{2}=1,\left\|X_{i j}\right\|_{\psi_{2}} \leq B$. Then for any $\varepsilon \in(0,1)$,

$$
\mathbb{P}\left(s_{n}(\Gamma) \leq \varepsilon n^{-1 / 2}\right) \leq C \varepsilon+c^{n}
$$

where $C>0, c \in(0,1)$ depend only on $B$.

## Theorem (Guédon, Litvak, Pajor, Tomczak-Jaegermann, R.A.)

Let $\Gamma$ be an $n \times n$ random matrix with independent isotropic $\log$-concave rows. Then, for any $\varepsilon \in(0,1)$,

$$
\mathbb{P}\left(s_{n}(\Gamma) \leq \varepsilon n^{-1 / 2}\right) \leq C \varepsilon+C \exp \left(-c n^{c}\right)
$$

and

$$
\mathbb{P}\left(s_{n}(\Gamma) \leq \varepsilon n^{-1 / 2}\right) \leq C \varepsilon^{n /(n+2)} \log ^{C}(2 / \varepsilon) .
$$

## Corollary

For any $\delta \in(0,1)$ there exists $C_{\delta}$ such that for any $n$ and $\varepsilon \in(0,1)$,

$$
\mathbb{P}\left(s_{n}(\Gamma) \leq \varepsilon n^{-1 / 2}\right) \leq C_{\delta} \varepsilon^{1-\delta}
$$

## Definition

For an $n \times n$ matrix $\Gamma$ define the condition number $\kappa(\Gamma)$ as

$$
\kappa(\Gamma)=\|\Gamma\| \cdot\left\|\Gamma^{-1}\right\|=\frac{s_{1}(\Gamma)}{s_{n}(\Gamma)} .
$$

## Corollary

If $\Gamma$ has independent isotropic log-concave columns, then for any $\delta>0, t>0$,

$$
\mathbb{P}(\kappa(\Gamma) \geq n t) \leq \frac{C_{\delta}}{t^{1-\delta}}
$$

## Thank you

